

LAB 2 18/03/2022
MIPS Sample Programs

1. MIPS assembly, write an assembly language version of the following C code segment

```
int A[100], B[100];
for (i=1; i < 100; i++) {
    A[i] = A[i-1] + B[i];
}
```

At the beginning of this code segment, the only values in registers are the base address of array A and B in registers \$a0 and \$a1. Avoid the use of multiplication instructions-they are unnecessary.

The MIPS assembly sequence is as follows

```
li $t0, 1      # Starting index of i
li $t5, 100    # Loop bound

loop:
lw $t1, 0($a1) # Load A[i-1]
lw $t2, 4($a2)  # Load B[i]
add $t3, $t1, $t2 # A[i-1] + B[i]
sw $t3, 4($a1) # A[i] = A[i-1] + B[i]
addi $a1, 4     # Go to i+1
addi $a2, 4     # Go to i+1
addi $t0, 1      # Increment index variable
bne $t0, $t5, loop # Compare with Loop Bound

halt:
nop
```

2. Convert the following C statements to equivalent MIPS assembly language. Assume that the variables f, g, I and j are assigned to registers \$s0, \$s1, \$s2 and \$s3 respectively. Assume that the base address of the array A and B are in registers \$s6 and \$s7 respectively.

a) $f = g + h + B[4]$

```
lw $t0, 16($s7)
add $s0, $s1, $s2
add $s0, $s0, $t0
```

b) $f = g - A[B[4]]$

```
lw $t0, 16($s7)
sll $t1, $t0, 2
add $t2, $t1, $s6
lw $t3, 0($t2)
sub $s0, $s1, $t3
```

Convert the following program into machine code.

```
0xFC00000C      start: .....
0xFC000010      loop: addi $t0,$t0,-1
0xFC000014          sw $t0, 4($t2)
0xFC000018          bne $t0, $t3, loop
0xFC00001C          j start
```

addi \$t0,\$t0,-1
-1 = 0xFFFF

0010 00/01 000/0 1000 /1111 1111 1111 1111
 8 8 8 -1
Op(6) rs(5) rt(5) immediate(16)

sw \$t0, 4(\$t2)
1010 11/01 010/0 1000 /0000 0000 0000 0100
 43 10 8 4
Op(6) rs(5) rt(5) immediate(16)

bne \$t0, \$t3, loop
target address = (immediate * 4) + address of the following instruction
immediate = (target address – address of the following instruction) / 4
= (FC000010 – FC00001C) / 4
= - C / 4
= -3 → 1111 1111 1111 1101

Or convert to binary first

= 1111 1100 0000 0000 0000 0000 0001 0000 –
 1111 1100 0000 0000 0000 0000 0001 1100

= 1111 1100 0000 0000 0000 0000 0001 0000 +
 0000 0011 1111 1111 1111 1111 110 0100

= 1111 1111 1111 1111 1111 1111 1111 0100 / 4 → srl by 2
= 11 1111 1111 1111 1111 1111 1101

immediate is 16 only so immediate = 1111 1111 1111 1101

0001 01/01 000/0 1011 /1111 1111 1111 1101
 5 8 11 -3
Op(6) rs(5) rt(5) immediate(16)

j start
target address = last 4 bits of PC : (immediate * 4)
immediate = first 28 bits from target address / 4
= C00000C / 4 = 3000003
= 1100 0000 0000 0000 0000 0000 1100 / 4 → srl by 2
= 11 0000 0000 0000 0000 0000 0011

0000 10/11 0000 0000 0000 0000 0000 0011
 2 3000003
Op(6) immediate (26)

4 factorial Program

```
data
    msg: .asciiiz "Enter a number"
    answer: .asciiiz "\nFactorial is: "
.text
# welcome message
li $v0, 4
la $a0, msg
syscall
# read integer
li $v0, 5
syscall
# print the integer
move $a0, $v0
li $v0, 1
syscall
jal calculate_factorial
move $a1, $v0
li $v0, 4
la $a0, answer
syscall
move $a0, $a1
li $v0, 1
syscall
li $v0, 10
syscall
calculate_factorial:
    addi $sp, $sp-4
    sw $ra, ($sp)
    li $v0, 1
    multiply:
        beq $a0, $zero, return
        mul $v0, $v0, $a0
        addi $a0, $a0, -1
        j multiply
    return:
    lw $ra, ($sp)
    jr $ra
```